Cryptography and Network Security Chapter 14

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#### Chapter 14 – Key Management and Distribution

No Singhalese, whether man or woman, would venture out of the house without a bunch of keys in his hand, for without such a talisman he would fear that some devil might take advantage of his weak state to slip into his body. —The Golden Bough, Sir James George

Frazer

# Key Management and Distribution

> topics of cryptographic key management / key distribution are complex cryptographic, protocol, & management issues symmetric schemes require both parties to share a common secret key > public key schemes require parties to acquire valid public keys >have concerns with doing both

### **Key Distribution**

>symmetric schemes require both parties to share a common secret key >issue is how to securely distribute this key whilst protecting it from others >frequent key changes can be desirable > often secure system failure due to a break in the key distribution scheme

# **Key Distribution**

- Given parties A and B have various **key distribution** alternatives:
  - A can select key and physically deliver to *B*.
  - A third party can select the key and physically deliver it to *A* and *B*.
  - If *A* and *B* have previously and recently used a key, one party can transmit the new key to the other, encrypted using the old key.
  - If *A* and *B* each has an encrypted connection to a third party *C*, *C* can deliver a key on the encrypted links to *A* and *B*.

# **Key Distribution Task**

- If end-to-end encryption is done at a network or IP level, then a key is needed for each pair of hosts.
- If there are N hosts, the number of required keys is [N(N-1)/2].



# **Key Hierarchy**

>typically have a hierarchy of keys

#### session key

- temporary key
- used for encryption of data between users
- for one logical session then discarded

#### master key

- used to encrypt session keys
- shared by user & key distribution center

# **Key Hierarchy**

The use of a key distribution center is based on the use of a hierarchy of keys.

➤ session key

temporary key

used for encryption of data between users for one logical session then discarded

master key

- used to encrypt session keys
- shared by user & key distribution centre

### Key Hierarchy



# **Key Distribution Scenario**

The scenario assumes that each user shares a unique master key with the key distribution center (KDC).



### **Key Distribution Issues**

- Hierarchies of KDC's : required for large networks, but must trust each other
  Session key lifetimes: should be limited for greater security
- use of automatic key distribution on behalf of users, but must trust system
- use of decentralized key distribution
- >controlling key usage

#### A Transparent Key Control Scheme



#### **Benefits**

The automated key distribution approach provides the flexibility and dynamic characteristics needed to allow a number of terminal users to access a number of hosts and for the hosts to exchange data with each other.

#### decentralized key distribution



- each end system be able to communicate in a secure manner with all potential partner end systems for purposes of session key distribution.
- each node must maintain at most (n 1) master keys, as many session keys as required may be generated and used

# Symmetric Key Distribution Using Public Keys

> public key cryptosystems are inefficient

- so almost never use for direct data encryption
- rather use to encrypt secret keys for distribution

### **Simple Secret Key Distribution**

Merkle proposed this very simple scheme

- allows secure communications
- no keys before/after exist



#### Man-in-the-Middle Attack

This very simple scheme is vulnerable to an active man-in-the-middle attack



#### Secret Key Distribution with Confidentiality and Authentication



# **Hybrid Key Distribution**

- retain use of private-key KDC
- >shares secret master key with each user
- >distributes session key using master key
- >public-key used to distribute master keys
  - especially useful with widely distributed users
- rationale
  - performance
  - backward compatibility

### **Distribution of Public Keys**

can be considered as using one of:

- public announcement
- publicly available directory
- public-key authority
- public-key certificates

#### **Public Announcement**

- users distribute public keys to recipients or broadcast to community at large
  - eg. append PGP keys to email messages or post to news groups or email list
- major weakness is forgery
  - anyone can create a key claiming to be someone else and broadcast it
  - until forgery is discovered can masquerade as claimed user

#### **Publicly Available Directory**

can obtain greater security by registering keys with a public directory

>directory must be trusted with properties:

- <u>contains</u> {name,public-key} entries
- participants register securely with directory
- participants <u>can replace key</u> at any time
- directory is periodically published
- directory can be <u>accessed electronically</u>

still vulnerable to tampering or forgery

## **Public-Key Authority**

- improve security by tightening control over distribution of keys from directory
- has properties of directory
- > and requires users to know public key for the directory
- then users interact with directory to obtain any desired public key securely
  - does require real-time access to directory when keys are needed
  - may be vulnerable to tampering



- 1. A sends a timestamped message to the public-key authority containing a request for the current public key of B.
- 2. The authority responds with a message that is encrypted using the authority's private key,  $PR_{\text{auth}}$ .
  - B's public key, *PU*<sub>b</sub>, which A can use to encrypt messages destined for B
  - The original request used to enable A to match this response with the corresponding earlier request and to verify that the original request was not altered before reception by the authority
  - The original timestamp given so A can determine that this is not an old message from the authority containing a key other than B's current public key
- **3.** A stores B's public key and also uses it to encrypt a message to B containing an identifier of A  $(ID_A)$  and a nonce  $(N_1)$ , which is used to identify this transaction uniquely.

- 4, 5. B retrieves A's public key from the authority in the same manner as A retrieved B's public key.
- 6. B sends a message to A encrypted with PU<sub>a</sub> and containing A's nonce (N<sub>1</sub>) as well as a new nonce generated by B (N<sub>2</sub>). Because only B could have decrypted message (3), the presence of N<sub>1</sub> in message (6) assures A that the correspondent is B.
- 7. A returns N<sub>2</sub>, which is encrypted using B's public key, to assure B that its correspondent is A.

### **Public-Key Certificates**

certificates allow key exchange without real-time access to public-key authority > a certificate binds identity to public key usually with other info such as period of validity, rights of use etc > with all contents **signed** by a trusted Public-Key or Certificate Authority (CA) can be verified by anyone who knows the public-key authorities public-key



# X.509 Authenticatio

Each certificate contains the public key of a user and is signed with the private key of a trusted certification authority.

- part of CCITT X.500 that defines service standards
  - distributed servers maintaining user info database
- > defines framework for authentication services
  - directory may store public-key certificates
  - with public key of user signed by certification authority
- > also defines authentication protocols
- uses public-key crypto & digital signatures
  - algorithms not standardised, but RSA recommended
- X.509 certificates are widely used
  - have 3 versions

# X.509 Certificate Use



signature using public key.

### X.509 Certificate

- issued by a Certification Authority (CA), co
  - version V (1, 2, or 3)
  - serial number <u>SN</u> (unique within CA) identifying
  - signature algorithm identifier <u>AI</u>
  - issuer X.500 name (CA)
  - period of validity <u>TA</u> (from to dates)
  - subject X.500 name A (name of owner)
  - subject public-key info Ap (algorithm, parameters, key) Ap
  - issuer unique identifier (v2+) UCA
  - subject unique identifier (v2+) UA
  - extension fields (v3)
  - signature (of hash of all fields in certificate)

notation CA<<A>> denotes certificate for A signed by CA

#### CA{V,SN,AI,CA,UCA,A,UA,Ap,TA}

A certificate contains a public key. The certificate, in addition to the pukey, contains additional information as issuer, what it's supposed to be for, and any other type of metadata. Typically a certificate is itself signed a private key, that verifies its authen

#### **X.509 Certificates**





(b) Certificate Revocation List

(a) X.509 Certificate

#### **Obtaining a Certificate**

any user with access to CA can get any certificate from it

- >only the CA can modify a certificate
- because cannot be forged, certificates can be placed in a public directory

#### if both users share a common CA then they are assumed to know its public key

#### > otherwise CA's must form a hierarchy

- use certificates linking members of hierarchy to validate other CA's
  - each CA has certificates for clients (forward) and parent (backward)

#### > each client trusts parents certificates

enable verification of any certificate from one CA by users of all other CAs in hierarchy

#### **CA Hierarchy Use**





In this example, user A can acquire the following certificates from the directory to establish a certification path to B:

#### $X \ll W \gg W \ll V \gg V \ll Y \gg Y \ll Z \gg Z \ll B \gg$

B require A's public key which can be obtained from the following certification path:

 $Z \ll Y \gg Y \ll V \gg V \ll W \gg W \ll X \gg X \ll A \gg$ 

#### **Certificate Revocation**

- certificates have a period of validity
- may need to revoke before expiry, eg:
  - user's private key is compromised
  - 2. user is no longer certified by this CA
  - 3. CA's certificate is compromised
- CA's maintain list of revoked certificates
  - the Certificate Revocation List (CRL)
- users should check certificates with CA's CRL

### X.509 Version 3

has been recognised that additional information is needed in a certificate email/URL, policy details, usage constraints rather than explicitly naming new fields defined a general extension method > extensions consist of: extension identifier criticality indicator extension value

#### **Certificate Extensions**

key and policy information

- convey info about subject & issuer keys, plus indicators of certificate policy
- certificate subject and issuer attributes
  - support alternative names, in alternative formats for certificate subject and/or issuer

certificate path constraints

 allow constraints on use of certificates by other CA's

### **Public Key Infrastructure**



# **PKIX Management**

#### functions:

- registration
- initialization
- certification
- key pair recovery
- key pair update
- revocation request
- cross certification

>protocols: CMP, CMC

# Summary

have considered:

- symmetric key distribution using symmetric encryption
- symmetric key distribution using public-key encryption
- distribution of public keys
  - announcement, directory, authrority, CA
- X.509 authentication and certificates
- public key infrastructure (PKIX)